

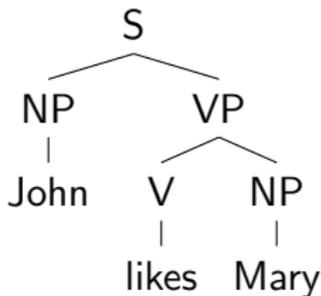
Categorical Grammar

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Phrase Structure Grammar



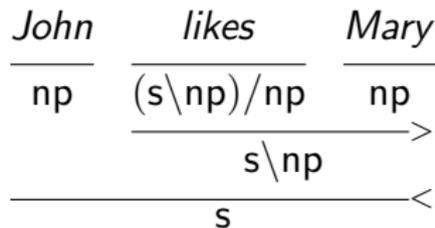
Last two lectures

- type theory
- argument structure
- λ calculus

Today's lecture: Categorical Grammar

$$\begin{array}{c}
 \frac{John}{np} \quad \frac{likes}{(s \backslash np) / np} \quad \frac{Mary}{np} \\
 \hline
 \frac{ \quad \quad }{s \backslash np} > \\
 \hline
 s <
 \end{array}$$

What is the meaning of the symbols in this analysis?



Complex category

In a CG, all constituents—and in particular the lexical elements—are associated with a **very specific category** which define their syntactic behaviour.

Simple *phrase-structure* rules

A set of **universal rules** defines how words and other constituents can be combined **according to their categories**.

Syntax vs. Semantics

Syntactic and **semantic** descriptions are tightly connected → CG is popular amongst logicians and semanticists.

Definition

The set of syntactic categories \mathcal{C} is defined recursively:

- **Atomic categories:** the grammar for each language is assumed to define a finite set of atomic categories, usually $s, np, n, pp, \dots \in \mathcal{C}$
- **Complex categories:** if X and $Y \in \mathcal{C}$, then $X/Y, X \setminus Y \in \mathcal{C}$

Complex categories X/Y or $X \setminus Y$ are functors

- X : a result
- Y : an argument
- $/$: arguments to the right of the functor
- \setminus : arguments to the left of the functor

Complex categories encode subcategorisation information

- intransitive verb: $s \backslash np$ ▷ walked
- transitive verb: $(s \backslash np) / np$ ▷ respected
- ditransitive verb: $((s \backslash np) / np) / np$ ▷ gave

$(s \backslash np) / np$

- the verb takes a noun phrase to its right, and
- another noun phrase to its left to form a sentence.

There is no explicit difference made between phrases and words:
An **intransitive verb** is described in the same way as a **verb phrase with an object**: $s \backslash np$.

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Modification

In CG, adjuncts have the following general form: $X \backslash X$ or X / X .

Example

- PP nominal: $(np \backslash np) / np$
- PP verbal: $((s \backslash np) \backslash (s \backslash np)) / np$

Lexicalization

In a lexicalized grammar, each element of the grammar contains at least one lexical item (terminal).

- G1: $S \rightarrow SS$, $S \rightarrow a$
- G2: $S \rightarrow aS$, $S \rightarrow a$

Grammar or Lexicon

In a CG,

- the lexicon specifies the categories that the words of a language can take;
- lexical entries do most of the grammatical work of mapping the strings of the language to their interpretations.

The Principle of Lexical Head Government

Both bounded and unbounded syntactic dependencies are specified by the lexical syntactic type of their head.

Example

- (1) a. *John* \vdash np ▷ *John* is a noun phrase.
 b. *shares* \vdash np ▷ *shares* is a noun phrase.
 c. *buys* \vdash (s\np)/np ▷ *buy* is a transitive verb.
 d. *sleeps* \vdash s\np ▷ *sleeps* is an intransitive verb.
 e. *well* \vdash (s\np)/(s\np) ▷ *well* can modify a
s\np-like thing.

The Principle of Head Categorial Uniqueness

A single nondisjunctive lexical category for the head of a given construction specifies both the bounded dependencies that arise when its complements are in canonical position and the unbounded dependencies that arise when those complements are displaced under relativization, coordination, and the like.

admire ⊢ (s\np)/np

- (2) a. John **admires** **Mary**.
 b. **the man** that I believe that John **admires**.
 c. I believe that John **admires** and you believe that he dislikes, **the woman in the skinny skirt**.

Variants of categorial grammar differ in the rules they allow.

- The system defined by **Ajdukiewicz** (1935) and **Bar-Hillel** (1953) forms the basis for all variants of categorial grammar.
- In AB categorial grammar, categories can only combine through function application.

Forward application

$$X/Y \quad Y \Rightarrow X \quad (>)$$

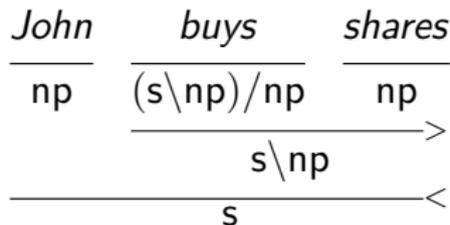
Backward application

$$Y \quad X \backslash Y \Rightarrow X \quad (<)$$

Deriving a string

- A string α is grammatical if each word in the string can be assigned a category (as defined by the **lexicon**) so that the lexical categories of the words in α can be combined (according to the **grammar rules**) to form a constituent.
- The process of combining constituents in this manner is called a derivation.

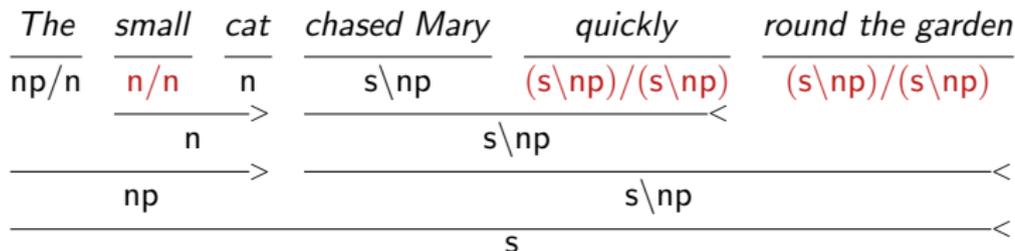
Example



Modification

In CG, adjuncts have the following general form: $X \backslash X$ or X / X .

Example



Not all $X \backslash X$ nor X / X are modifiers.

L. Bloomfield, *Language*

The lexicon is really an appendix of the grammar, a list of basic irregularities. This is all the more evident if meanings are taken into consideration, since the meaning of each morpheme belongs to it by an arbitrary tradition.

CG's view

- If this is the case, nothing in the lexicon is predictable, hence we do not need a theory of the lexicon.
- CG argues that this dichotomy gets in the way of our understanding of how syntax can shape possible lexicons.
 - * Any combinatory difference must be lexically specifiable.

Architecture

Syntactic structure (CG derivation) + Lexical interpretation



Meaning representation

$buy \vdash (s \setminus np_1) / np_2 : buy \rightarrow_A np_1 \wedge buy \rightarrow_P np_2$

- Syntactic category: $(s \setminus np) / np$
- Semantics (argument structure): the np indexed with “2” is the Patient of *buy*; the np indexed with “1” is the Agent of *buy*.

Semantic interpretation (2)

$buy \vdash (s \setminus np_1) / np_2 : buy \rightarrow_A np_1 \wedge buy \rightarrow_P np_2$

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Using a Dependency Interpretation

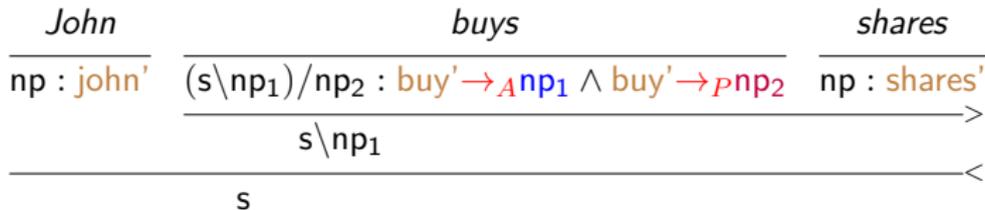


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Using a Dependency Interpretation

$John$	$buys$	$shares$
$np : john'$	$(s \setminus np_1) / np_2 : buy' \rightarrow_A np_1 \wedge buy' \rightarrow_P np_2$	$np : shares'$
$s \setminus np_1 : buy' \rightarrow_A np_1 \wedge buy' \rightarrow_P shares'$		
s		

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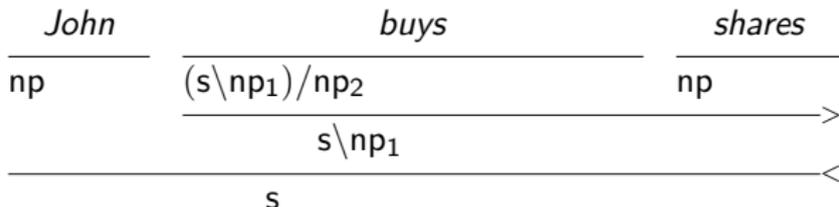
Using a Dependency Interpretation

$$\frac{\frac{John}{np : john'} \quad \frac{buys}{(s \setminus np_1) / np_2 : buy' \rightarrow_A np_1 \wedge buy' \rightarrow_P np_2} \quad \frac{shares}{np : shares'}}{s \setminus np_1 : buy' \rightarrow_A np_1 \wedge buy' \rightarrow_P shares'} >$$
$$\frac{}{s : buy' \rightarrow_A john' \wedge buy' \rightarrow_P shares'} <$$

$buy \vdash (s \setminus np_1) / np_2 : \lambda x. \lambda y. buy'(y, x)$

- Syntactic category: $(s \setminus np) / np$
- Semantics (λ expression): the np indexed with “2” is the second argument of buy ; the np indexed with “1” is the first argument of buy .

Using λ calculus



Semantic interpretation (3)

$\text{buy} \vdash (\text{s} \backslash \text{np}_1) / \text{np}_2 : \lambda x. \lambda y. \text{buy}'(y, x)$

- Syntactic category: $(\text{s} \backslash \text{np}) / \text{np}$
- Semantics (λ expression): the **np** indexed with “2” is the second argument of *buy*; the **np** indexed with “1” is the first argument of *buy*.

Using λ calculus

$$\frac{\frac{\text{John}}{\text{np} : \text{john}'}}{\text{np} : \text{john}'}}{\frac{\frac{\text{buys}}{(\text{s} \backslash \text{np}_1) / \text{np}_2 : \lambda x. \lambda y. \text{buy}'(y, x)}}{\text{s} \backslash \text{np}_1}}{\frac{\frac{\text{shares}}{\text{np} : \text{shares}'}}{\text{np} : \text{shares}'}}{\text{s}}}$$

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$$\frac{\frac{\text{John}}{\text{np} : \text{john}'}}{\frac{\frac{\text{buys}}{(\text{s} \backslash \text{np}_1) / \text{np}_2 : \lambda x. \lambda y. \text{buy}'(y, x)}}{\text{np} : \text{shares}'}} \rightarrow}{\text{s} \backslash \text{np}_1 : \lambda y. \text{buy}'(y, \text{shares}')} \leftarrow$$

s

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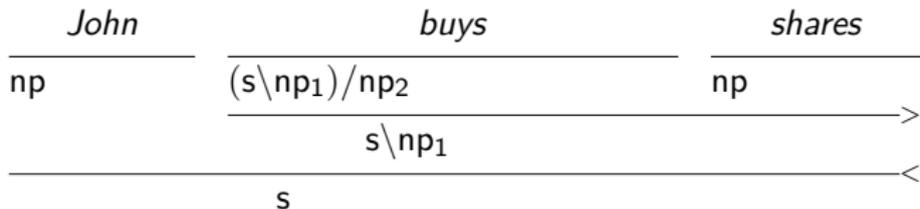
Using λ calculus

$$\frac{\frac{John}{np : john'} \quad \frac{buys \quad \frac{shares}{np : shares'}}{(s \setminus np_1) / np_2 : \lambda x. \lambda y. buy'(y, x)}}{s \setminus np_1 : \lambda y. buy'(y, shares')} \rightarrow}{s : buy'(john', shares')} \leftarrow$$

$\text{buy} \vdash (\text{s} \setminus \text{np}_1) / \text{np}_2 : \lambda x. \lambda y. \text{buy}'(y, x)$

- Syntactic category: $(\text{s} \setminus \text{np}) / \text{np}$
- Semantic type (type theory): “buy” $\Rightarrow \langle e, \langle e, t \rangle \rangle$
 Type theory could easily admit and handle n -place functors;
 such a functor can take the n arguments one by one

Using types

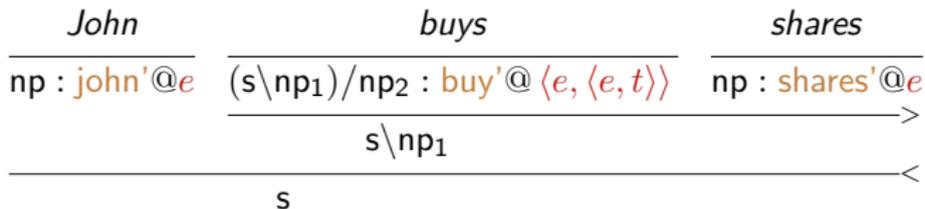


Semantic interpretation (4)

$\text{buy} \vdash (\text{s} \setminus \text{np}_1) / \text{np}_2 : \lambda x. \lambda y. \text{buy}'(y, x)$

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Using types



Semantic interpretation (4)

$\text{buy} \vdash (\text{s} \setminus \text{np}_1) / \text{np}_2 : \lambda x. \lambda y. \text{buy}'(y, x)$

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Type theory could easily admit and handle n -place functors;
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Using types

$$\frac{\frac{\text{John}}{\text{np} : \text{john}'@e} \quad \frac{\text{buys}}{(\text{s} \setminus \text{np}_1) / \text{np}_2 : \text{buy}'@ \langle e, \langle e, t \rangle \rangle} \quad \frac{\text{shares}}{\text{np} : \text{shares}'@e}}{\text{s} \setminus \text{np}_1 : \text{buy}'(\text{shares}')@ \langle e, t \rangle} \begin{array}{l} \longrightarrow \\ \longleftarrow \end{array}$$

$\text{buy} \vdash (\text{s} \backslash \text{np}_1) / \text{np}_2 : \lambda x. \lambda y. \text{buy}'(y, x)$

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 \hline
 \text{s} \backslash \text{np}_1 : \text{buy}'(\text{shares}')@ \langle e, t \rangle \quad \text{>} \\
 \hline
 \text{s} : \text{buy}'(\text{john}')(\text{shares}')@t \quad \text{<}
 \end{array}$$

The principle of type transparency

The principle of Categorical Type Transparency

For a given language, **the semantic type** of the interpretation together with a number of language-specific directional parameter settings uniquely determines **the syntactic category** of a category.

The inverse of Type Transparency

For any category, the semantic type is a function of the syntactic type.

- CGs put into the lexicon most of the information that is captured in CFG rules.
- In CGs, all constituents and **lexical elements** are associated with a syntactic “**category**.”
- In CGs, syntactic information is tightly related to semantic information.

Example

S → NP VP

VP → TV NP

TV → married|finds|...

married := (s\np)/np.